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# On the minimum of the unit lattice. 

par Volker KESSLER

## 1. Introduction.

Computations in lattices often require a lower bound for the minimum of the lattice, both for practical purposes and for a theoretical analysis of the algorithms, e.g. [1] and [2].

In this paper we recall two results of Dobrowolski [3] and Smyth [5] in order to get such a bound for the unit lattice.

## 2. Lower bound.

Let $K$ be a finite extension of $\mathbb{Q}$ of degree $n$ with maximal order $R$. For $1 \leq i \leq n$ we denote by

$$
K \rightarrow K^{(i)} \subset \mathbb{C}, \quad \alpha \rightarrow \alpha^{(i)}
$$

the $n$ different embeddings of $K$ into the field $\mathbb{C}$ of complex numbers. The first $r_{1}$ of those embeddings are real, the last $2 r_{2}$ embeddings are non-real and numbered such that the $\left(r_{1}+r_{2}+i\right)$ th embedding is the complexconjugation of the $\left(r_{1}+i\right)$ th embedding. Then the logarithmic map is given by

$$
\log : K^{*} \rightarrow \mathbb{R}^{r}, \quad \log (\alpha):=\left(c_{1} \log \left|\alpha^{(1)}\right|, \cdots, c_{r} \log \left|\alpha^{(r)}\right|\right)
$$

with the unit rank $r=r_{1}+r_{2}-1$ and

$$
c_{i}= \begin{cases}1 & \text { for } 1 \leq i \leq r_{1} \\ 2 & \text { for } r_{1}+1 \leq i \leq r+1\end{cases}
$$

The kernel of Log consists exactly of the roots of the unity lying in $K$. We define the minimum $\lambda(L)$ of the unit lattice $L:=\log \left(R^{*}\right)$ by

$$
\lambda(L)=\min \{\|v\| \mid v \in L \backslash\{0\}\}
$$

[^0]where || || denotes the Euclidean norm.
Theorem : A lower bound for the minimum $\lambda(L)$ is given by (1)
$$
\lambda(L)>\mu(K):=\sqrt{\frac{2}{r+1}}\left(\frac{1}{1200}\left(\frac{\log \log n}{\log n}\right)^{3}-\frac{1}{2880000}\left(\frac{\log \log n}{\log n}\right)^{6}\right)
$$
which is "a bit" larger than
$$
\frac{1}{\sqrt{r+1}} \frac{1}{1000}\left(\frac{\log \log n}{\log n}\right)^{3}
$$

Thus the inverse $1 / \lambda(L)$ is of the magnitude $0\left(n^{1 / 2+¢}\right)$ for every $\epsilon>0$.
Proof. Let $\epsilon \in R^{*}$ be a unit of degree $m$ over $\mathbb{Q}$, which is no root of unity. Without loss of generality we can assume that $m=n$, because if $\|\log \epsilon\|$ is larger than $\mu\left(K^{\prime}\right)$ for a subfield $K^{\prime}$ of $K$ it is also larger than $\mu(K)$.

We are interested in two subsets of the conjugates $\epsilon^{(1)}, \cdots, \epsilon^{(n)}$

$$
\begin{aligned}
& S:=\left\{1 \leq i \leq r+1| | \epsilon^{(i)} \mid>1\right\} \\
& T:=\left\{1 \leq i \leq r+1| | \epsilon^{(i)} \mid<1\right\}
\end{aligned}
$$

Since $\epsilon$ is no root of unity $S$ is non-empty and therefore $T$ cannot be empty because of $N(\epsilon)=1$.

We call $\epsilon$ reciprocal if $\epsilon$ is conjugate to $\epsilon^{-1}$, i.e. its minimal polynomial $f(X)=X^{n}+a_{n-1} X^{n-1}+\cdots+a_{0}$ satisfies

$$
f(X)=X^{n} f\left(\frac{1}{X}\right)=a_{0} X^{n}+a_{1} X^{n-1}+\cdots+a_{n-1} X+1
$$

If $\epsilon$ is non-reciprocal we know from the theorem of [5] that

$$
\prod_{i \in S}\left|\epsilon^{(i)}\right|^{c_{i}} \geq \theta
$$

where $\theta$ is the real root of $X^{3}-X-1$, i.e. $\theta \approx 1.3247$. Thus

$$
\begin{equation*}
\sum_{i \in S} c_{i} \log \left|\epsilon^{(i)}\right| \geq \log \theta \approx 0.281 \tag{2}
\end{equation*}
$$

But from $N(\epsilon)=1$ it follows

$$
\begin{equation*}
\sum_{i \in S} c_{i} \log \left|\epsilon^{(i)}\right|=-\sum_{i \in T} c_{i} \log \left|\epsilon^{(i)}\right| \tag{3}
\end{equation*}
$$

The value $c_{r+1} \log \left|\epsilon^{(r+1)}\right|$ does not occur in the norm of $\log (\epsilon)$. But as a consequence of (3) it does not matter if $r+1$ lies in $S$ or in $T$ and so we can assume without restriction that $r+1 \notin S$. Thus

$$
\begin{aligned}
\|\log (\epsilon)\| & \geq \sqrt{\sum_{i \in S}\left(c_{i} \log \left|\epsilon^{(i)}\right|\right)^{2}} \\
& \geq r^{-1 / 2} \sum_{i \in S}\left(c_{i} \log \left|\epsilon^{(i)}\right|\right) \geq r^{-1 / 2} \log \theta>\mu(K)
\end{aligned}
$$

(The second inequality follows from the well known norm equivalence between 1-norm and Euclidean norm.)

For reciprocal $\epsilon$ we know by Theorem 1 of [3]:

$$
\begin{equation*}
\prod_{i \in S}\left|\epsilon^{(i)}\right|^{c_{i}}>1+\frac{1}{1200}\left(\frac{\log \log n}{\log n}\right)^{3} \tag{4}
\end{equation*}
$$

We now use the Taylor series of the logarithm $(|y|<1)$ :

$$
\begin{equation*}
\log (1+y)=y-\frac{y^{2}}{2}+\frac{y^{3}}{3} \mp \cdots>y-\frac{y^{2}}{2} \tag{5}
\end{equation*}
$$

The inequality follows directly from Lagrange's representation of the residue. Applying (5) to (4) yields

$$
\sum_{i \in S} c_{i} \log \left|\epsilon^{(i)}\right|>\frac{1}{1200}\left(\frac{\log \log n}{\log n}\right)^{3}-\frac{1}{2880000}\left(\frac{\log \log n}{\log n}\right)^{6}
$$

Since $\epsilon$ is reciprocal the inverses of the conjugates of $\epsilon$ are also conjugate to $\epsilon$. This implies that the numbers of conjugates outside the unit circle equals the number of conjugates inside the unit circle, i.e

$$
\# S=\# T \leq \frac{r+1}{2} \leq \frac{n}{2}
$$

Again by (3) we can assume that $r+1 \notin S$

$$
\begin{aligned}
& \|\log (\epsilon)\| \geq \sqrt{\sum_{i \in S}\left(c_{i} \log \left|\epsilon^{(i)}\right|\right)^{2}} \geq \sqrt{\frac{2}{r+1}} \sum_{i \in S} c_{i} \log \left|\epsilon^{(i)}\right| \\
& >\sqrt{\frac{2}{r+1}}\left(\frac{1}{1200}\left(\frac{\log \log n}{\log n}\right)^{3}-\frac{1}{2880000}\left(\frac{\log \log n}{\log n}\right)^{6}\right)=\mu(K)
\end{aligned}
$$

which is larger than

$$
\sqrt{\frac{2}{r+1}}\left(\frac{1}{1200}-\frac{1}{2880000}\right)\left(\frac{\log \log n}{\log n}\right)^{3} .
$$

Because of $\sqrt{2}\left(\frac{1}{1200}-\frac{1}{2880000}\right) \approx 0.001178$ we thus proved the lower bound.
Remark. If the conjecture of Schinzel and Zassenhaus [5] is correct the term $\left(\frac{\log \log n}{\log n}\right)^{3}$ can be substituted by a constant independent of $n$. This bound would be provable the best one (up to constants).

## References

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